

# A non-commutative topology on $\text{rep } A$

Lieven Le Bruyn

Departement Wiskunde en Informatica

Universiteit Antwerpen

B-2020 Antwerp (Belgium)

lieven.lebruyn@ua.ac.be

## Abstract

We extend the Zariski topology on  $\text{simp } A$ , the set of all simple finite dimensional representations of  $A$ , to a non-commutative topology (in the sense of Fred Van Oystaeyen) on  $\text{rep } A$ , the set of all finite dimensional representations of  $A$ , using Jordan-Hölder filtrations. The non-commutativity of the topology is enforced by the order of the composition factors.

All algebras will be affine associative  $\mathbb{k}$ -algebras with unit over an algebraically closed field  $\mathbb{k}$ . The *non-commutative affine 'scheme'* associated to an algebra  $A$  is, as a set, the disjoint union

$$\text{rep } A = \bigsqcup_n \text{rep}_n A$$

where  $\text{rep}_n A$  is the (commutative) affine scheme of  $n$ -dimensional representations of  $A$ . In this note we will equip  $\text{rep } A$  with a non-commutative topology in the sense of Fred Van Oystaeyen [5, §7.2] (or, more precisely, a slight generalization of it).

Here is the main idea. The twosided prime ideal spectrum  $\text{spec } A$  is an (ordinary) topological space via the Zariski topology, see for example [4] or [1, §II.6]. Hence, the subset  $\text{simp } A$  of all simple finite dimensional  $A$ -representations can be equipped with the induced topology. This topology can then be extended to a non-commutative topology on  $\text{rep } A$  using Jordan-Hölder filtrations. The non-commutative nature of the topology is enforced by the order of the composition factors.

We give a few examples, connect this notion with that of Reineke's composition monoid and remark on the difference between quotient varieties and moduli spaces from the perspective of non-commutative topology. Finally, we note that this construction can be generalized verbatim to any Artinian Abelian category as soon as we have a topology on the set of simple objects.

## 1 The Zariski topology on $\text{simp } A$ .

Recall that a prime ideal  $P$  of  $A$  is a two-sided ideal satisfying the property that if  $I \cdot J \subset P$  then  $I \subset P$  or  $J \subset P$  for any pair of two-sided ideals  $I, J$  of  $A$ . The *prime spectrum*  $\text{spec } A$  is the set of all two-sided prime ideals of  $A$ . The *Zariski topology* on  $\text{spec } A$  has as its closed subsets

$$\mathbb{V}(S) = \{P \in \text{spec } A \mid S \subset P\}$$

where  $S$  varies over all subsets of  $A$ , see for example [1, Prop. II.6.2]. Note that an algebra morphism  $\phi : A \longrightarrow B$  does *not* necessarily induce a continuous map  $\phi^* : \text{spec } B \longrightarrow \text{spec } A$  but it does so in the case  $\phi$  is a *central extension* in the sense of [1, §II.6].

If  $M \in \text{rep}_n A$  is a simple  $n$ -dimensional representation, there is a defining epimorphism  $\psi_M : A \twoheadrightarrow M_n(\mathbb{k})$  and the kernel of this morphism  $\ker \psi_M$  is a two-sided maximal (hence prime) ideal of  $A$ . We define the Zariski topology on the set of all simple finite dimensional representations  $\text{simp } A$  by taking as its closed subsets

$$\mathbb{V}(S) = \{M \in \text{simp } A \mid S \subset \ker \psi_M\}$$

Again, one should be careful that whereas an algebra map  $\phi : A \longrightarrow B$  induces a map  $\phi^* : \text{rep } B \longrightarrow \text{rep } A$  it does *not* in general map  $\text{simp } B$  to  $\text{simp } A$  (unless  $\phi$  is a central extension).

With  $\mathcal{L}_A$  we will denote the set of all open subsets of  $\text{simp } A$ .  $\mathcal{L}_A$  will be the set of *letters* on which to base our non-commutative topology.

## 2 Non-commutative topologies (and generalizations).

In [5, Chp. 7] Fred Van Oystaeyen defined *non-commutative topologies* which are generalizations of usual topologies in which it is no longer true that  $A \cap A$  is equal to  $A$  for an open set  $A$ . In order to keep dichotomies of possible definitions to a minimum he imposed left-right symmetric conditions on the definition. However, for applications to representation theory it seems that the most natural non-commutative topologies are truly one-sided. For this reason we take some time to generalize some definitions and results of [5, Chp. 7].

We fix a partially ordered set  $(\Lambda, \leq)$  with a unique minimal element 0 and a unique maximal element 1, equipped with two operations  $\wedge$  and  $\vee$ . With  $i_\Lambda$  we will denote the set of all *idempotent elements* of  $\Lambda$ , that is, those  $x \in \Lambda$  such that  $x \wedge x = x$ . A *finite global cover* is a finite subset  $\{\lambda_1, \dots, \lambda_n\}$  such that  $1 = \lambda_1 \vee \dots \vee \lambda_n$ . In the table below we have listed the conditions for a (one-sided) non-commutative topology. Note that some requirements are less essential than others. For example, the covering condition (A10) is only needed if we want to fix non-commutative topologies in the framework of non-commutative Grothendieck topologies [5] and the weak modularity condition (A9) is not required if every basic open is  $\vee$ -idempotent (as is the case in most examples).

(A1)	$x \wedge y \leq x$		$x \wedge y \leq y$
(A2)	$x \wedge 1 = x$ $x \wedge 0 = 0$		$1 \wedge x = x$ $0 \wedge x = 0$
(A3)		$(x \wedge y) \wedge z = x \wedge (y \wedge z) = x \wedge y \wedge z$	
(A4)	$x \leq y \Rightarrow z \wedge x \leq z \wedge y$		$x \leq y \Rightarrow x \wedge z \leq y \wedge z$
(A5)	$x \leq x \vee y$		$y \leq x \vee y$
(A6)	$x \vee 1 = 1$ $x \vee 0 = x$		$1 \vee x = 1$ $0 \vee x = x$
(A7)		$(x \vee y) \vee z = x \vee (y \vee z) = x \vee y \vee z$	
(A8)	$x \leq y \Rightarrow x \vee z \leq y \vee z$		$x \leq y \Rightarrow z \vee x \leq z \vee y$
(A9)	$a \vee (a \wedge b) \leq (a \vee a) \wedge b$		$a \vee (b \wedge a) \leq (a \vee b) \wedge a$
(A10)	$x = (x \wedge \lambda_1) \vee \dots \vee (x \wedge \lambda_n)$		$x = (\lambda_1 \wedge x) \vee \dots \vee (\lambda_n \wedge x)$

**Definition 1** Let  $(\Lambda, \leq)$  be a partially ordered set with minimal and maximal element 0 and 1 and operations  $\wedge$  and  $\vee$ . Then,

$\Lambda$  is said to be a *left non-commutative topology* if and only if the left and middle column conditions of (A1)-(A10) are valid for all  $x, y, z \in \Lambda$ , all  $a, b \in i_\Lambda$  with  $a \leq b$  and all finite global covers  $\{\lambda_1, \dots, \lambda_n\}$ .

$\Lambda$  is said to be a *right non-commutative topology* if and only if the middle and right column conditions of (A1)-(A10) are valid for all  $x, y, z \in \Lambda$ , all  $a, b \in i_\Lambda$  with  $a \leq b$  and all finite global covers  $\{\lambda_1, \dots, \lambda_n\}$ .

$\Lambda$  is said to be a *non-commutative topology* if and only if the conditions (A1)-(A10) are valid for all  $x, y, z \in \Lambda$ , all  $a, b \in i_\Lambda$  with  $a \leq b$  and all finite global covers  $\{\lambda_1, \dots, \lambda_n\}$ .

There are at least two ways of building a genuine non-commutative topology out of these sets of basic opens. We briefly sketch the procedures here and refer to the forthcoming monograph [6] for details in the symmetric case (the one-sided versions present no real problems).

Let  $T(\Lambda)$  be the set of all finite  $(\wedge, \vee)$ -words in the *contractible* idempotent elements  $i_\Lambda$  (that is,  $\lambda \in i_\Lambda$  such that for all  $\lambda_1, \lambda_2$  with  $\lambda \leq \lambda_1 \vee \lambda_2$  we have that  $\lambda = (\lambda \wedge \lambda_1) \vee (\lambda \wedge \lambda_2)$ ). If  $\Lambda$  is a (left, right) non-commutative topology, then so is  $T(\Lambda)$ . The  $\vee$ -complete topology of virtual opens  $T'(\Lambda)$  is then the set of all  $(\wedge, \vee)$ -words in the contractible idempotents of finite length in  $\wedge$  (but not necessarily of finite length in  $\vee$ ). This non-commutative topology has properties very similar to that of an ordinary topology and, in fact, has associated to it a *commutative shadow*.

The second construction, leading to the *pattern topology*, starts with the equivalence classes of *directed systems*  $S \subset \Lambda$  (that is, if for all  $x, y \in S$  there is a  $z \in S$  such that  $z \leq x$  and  $z \leq y$ ) and where the equivalence relation  $S \sim S'$  is defined by

$$\begin{cases} \forall a \in S, \exists a' \in S, a' \leq a \text{ and } b \leq a' \leq b' \text{ for some } b, b' \in S' \\ \forall b \in S', \exists b' \in S', b' \leq b \text{ and } a \leq b' \leq a' \text{ for some } a, a' \in S \end{cases}$$

One can extend the  $\wedge, \vee$  operations on  $\Lambda$  to the equivalence classes  $C(\Lambda) = \{[S] \mid S \text{ directed}\}$  in the obvious way such that also  $C(\Lambda)$  is a (left, right) non-commutative topology. A directed set  $S \subset \Lambda$  is said to be *idempotent* if for all  $a \in S$ , there is an  $a' \in S \cap i_\Lambda$  such that  $a' \leq a$ . If  $S$  is idempotent then  $[S] \in i_{C(\Lambda)}$  and those idempotents will be called *strong idempotents*. The pattern topology  $\Pi(\Lambda)$  is the (left, right) non-commutative topology of finite  $(\wedge, \vee)$ -words in the strong idempotents of  $C(\Lambda)$ . A directed system  $[S]$  is called a *point* iff  $[S] \leq \vee[S_\alpha]$  implies that  $[S] \leq [S_\alpha]$  for some  $\alpha$ .

### 3 The basic opens.

For an  $n$ -dimensional representation  $M$  of  $A$  we call a finite filtration of length  $u$

$$\mathcal{F}^u : 0 = M_0 \subset M_1 \subset \dots \subset M_u = M$$

of  $A$ -representations a *Jordan-Hölder filtration* if the successive quotients

$$\mathcal{F}_i = \frac{M_i}{M_{i-1}}$$

are simple  $A$ -representations. Recall that  $\mathcal{L}_A$  is the set of all open subsets  $V$  of  $\text{simp } A$ . With  $\mathbb{W}_A$  we denote the non-commutative words in these letters

$$\mathbb{W}_A = \{V_1 \dots V_k \mid V_i \in \mathcal{L}_A, k \in \mathbb{N}\}$$

For a given word  $w = V_1 V_2 \dots V_k \in \mathbb{W}_A$  we define the *left basic open set*

$$\mathcal{O}_w^l = \{M \in \text{rep } A \mid \exists \mathcal{F}^u \text{ Jordan-Hölder filtration on } M \text{ such that } \mathcal{F}_i \in V_i\}$$

and the *right basic open set*

$$\mathcal{O}_w^r = \{M \in \text{rep } A \mid \exists \mathcal{F}^u \text{ Jordan-Hölder filtration on } M \text{ such that } \mathcal{F}_{u-i} \in V_{k-i}\}$$

Finally, to make these definitions symmetric we define the *basic open set*

$$\begin{aligned} \mathcal{O}_w = \{M \in \text{rep } A \mid \exists \mathcal{F}^u \text{ Jordan-Hölder filtration on } M \text{ such that } \mathcal{F}_{i_j} \in V_{j_j} \\ \text{for some } 1 \leq i_1 < i_2 < \dots < i_k \leq u \} \end{aligned}$$

Clearly,  $\mathcal{O}_w^l$  consists of those representations having prescribed bottom structure, whereas  $\mathcal{O}_w^r$  consists of those with prescribed top structure. In order to avoid three sets of definitions we will denote from now on  $\mathcal{O}_w^\bullet$  whenever we mean  $\bullet \in \{l, r, \emptyset\}$ .

If  $w = L_1 \dots L_k$  and  $w' = M_1 \dots M_l$ , we will denote with  $w \cup w'$  the *multi-set*  $\{N_1, \dots, N_m\}$  where each  $N_i$  is one of  $L_j, M_j$  and  $N_i$  occurs in  $w \cup w'$  as many times as its maximum number of factors in  $w$  or  $w'$ . With  $\text{rep}(w \cup w')$  we denote the subset of  $\text{rep } A$  consisting of the representations of  $M$  having a Jordan-Hölder filtration having factor-multi-set containing  $w \cup w'$ . For any triple of words  $w, w'$  and  $w''$  we denote  $\mathcal{O}_{w''}^\bullet(w \cup w') = \mathcal{O}_{w''}^\bullet \cap \text{rep}(w \cup w')$ .

We define an equivalence relation on the basic open sets by

$$\mathcal{O}_w^\bullet \approx \mathcal{O}_{w'}^\bullet \quad \Leftrightarrow \quad \mathcal{O}_w^\bullet(w \cup w') = \mathcal{O}_{w'}^\bullet(w \cup w')$$

The reason for this definition is that the condition of  $M \in \mathcal{O}_w^\bullet$  is void if  $M$  does not have enough Jordan-Hölder components to get all factors of  $w$  which makes it impossible to define equality of basic open sets defined by different words.

We can now define the partially ordered sets  $\Lambda_A^\bullet$  as consisting of all basic open subsets  $\mathcal{O}_w^\bullet$  of  $\text{rep } A$ . The partial ordering  $\leq$  is induced by set-theoretic inclusion modulo equivalence, that is,

$$\mathcal{O}_w^\bullet \leq \mathcal{O}_{w'}^\bullet \quad \Leftrightarrow \quad \mathcal{O}_w^\bullet(w \cup w') \subseteq \mathcal{O}_{w'}^\bullet(w \cup w')$$

As a consequence, equality  $=$  in the set  $\Lambda_A^\bullet$  coincides with equivalence  $\approx$ . Observe that these partially ordered sets have a unique minimal and a unique maximal element (upto equivalence)

$$0 = \emptyset = \mathcal{O}_\emptyset^\bullet \quad \text{and} \quad 1 = \text{rep } A = \mathcal{O}_{\text{simp } A}^\bullet$$

The operations  $\vee$  and  $\wedge$  are defined as follows :  $\vee$  is induced by ordinary set-theoretic union and  $\wedge$  is induced by concatenation of words, that is

$$\mathcal{O}_w^\bullet \wedge \mathcal{O}_{w'}^\bullet \approx \mathcal{O}_{ww'}^\bullet$$

**Theorem 1** *With notations as before :*

- $(\Lambda_A^l, \leq, \approx, 0, 1, \vee, \wedge)$  is a left non-commutative topology on  $\mathbf{rep} A$ .
- $(\Lambda_A^r, \leq, \approx, 0, 1, \vee, \wedge)$  is a right non-commutative topology on  $\mathbf{rep} A$ .

**Proof.** The tedious verification is left to the reader. Here, we only stress the importance of the equivalence relation for example in verifying  $x \wedge 1 = x$ . So, let  $w = L_1 \dots L_k$  then

$$\mathcal{O}_w^l \wedge 1 = \mathcal{O}_{L_1 \dots L_k \mathbf{simp} A}^l \subset \mathcal{O}_w^l$$

and this inclusion is proper (look at elements in  $\mathcal{O}_w^l$  having exactly  $k$  composition factors). However, as soon as the representation has  $k + 1$  composition factors, it is contained in the left hand side whence  $\mathcal{O}_w^l \wedge 1 \approx \mathcal{O}_w^l$ . A similar argument is needed in the covering condition.  $\square$

Note however that  $(\Lambda_A, \leq, \approx, 0, 1, \vee, \wedge)$  is not necessarily a non-commutative topology : the problematic conditions are  $\mathcal{O}_w \wedge 1 = \mathcal{O}_w = 1 \wedge \mathcal{O}_w$  and the covering condition. The reason is that for  $w = L_1 \dots L_k$  as before and  $M \in \mathcal{O}_w$  having  $> k$  factors, it may happen that the last factor is the one in  $L_k$  leaving no room for a successive factor in  $\mathbf{simp} A$  (whence  $\mathcal{O}_w \cap 1$  is not equivalent to  $\mathcal{O}_w$ ).

**Example 1** Let  $A$  be a finite dimensional algebra, then  $A$  has a finite number of simple representations  $\mathbf{simp} A = \{S_1, \dots, S_n\}$  and the Zariski topology is the discrete topology. If for some  $1 \leq i, j \leq n$  we have that

$$\mathrm{Ext}_A^1(S_i, S_j) = 0 \quad \text{and} \quad \mathrm{Ext}_A^1(S_j, S_i) \neq 0$$

then  $\Lambda_A^l$  is a genuinely non-commutative topology, for example

$$\mathcal{O}_{S_i}^l \wedge \mathcal{O}_{S_j}^l = \mathcal{O}_{S_i S_j}^l \neq \mathcal{O}_{S_j S_i}^l = \mathcal{O}_{S_j}^l \wedge \mathcal{O}_{S_i}^l$$

as a non-trivial extension  $0 \longrightarrow S_i \longrightarrow X \longrightarrow S_j \longrightarrow 0$  belongs to  $\mathcal{O}_{S_i S_j}^l(S_i S_j \cup S_j S_i)$  but not to  $\mathcal{O}_{S_j S_i}^l(S_i S_j \cup S_j S_i)$ .

## 4 Reineke's mon(str)oid.

When  $A$  is the path algebra of a quiver without oriented cycles we can generalize the foregoing example and connect the previous definitions to the *composition monoid* introduced and studied by Markus Reineke in [2].

Let  $Q$  be a quiver without oriented cycles, then its path algebra  $A = \mathbb{k}Q$  is finite dimensional hereditary with all simple representations one-dimensional and in one-to-one correspondence with the vertices of  $Q$ . For every dimension  $n$  we have that

$$\mathbf{rep}_n A = \bigsqcup_{|\alpha|=n} GL_n \times^{GL(\alpha)} \mathbf{rep}_\alpha Q$$

where  $\alpha$  runs over all dimension vectors of total dimension  $n$  and where  $\text{rep}_\alpha Q$  is the affine space of all  $\alpha$ -dimensional representations of the quiver  $Q$  with base-change group action by  $GL(\alpha)$ .

The *Reineke monstroid*  $\mathcal{M}(Q)$  has as its elements the set of all irreducible closed  $GL(\alpha)$ -stable subvarieties of  $\text{rep}_\alpha Q$  for all dimension vectors  $\alpha$ , equipped with a product

$$\mathcal{A} * \mathcal{B} = \{X \in \text{rep}_{\alpha+\beta} Q \mid \text{there is an exact sequence} \\ 0 \longrightarrow M \longrightarrow X \longrightarrow N \longrightarrow 0 \quad M \in \mathcal{A}, N \in \mathcal{B}\}$$

if  $\mathcal{A}$  (resp.  $\mathcal{B}$ ) is an element of  $\mathcal{M}(Q)$  contained in  $\text{rep}_\alpha Q$  (resp. in  $\text{rep}_\beta Q$ ). It is proved in [2, lemma 2.2] that  $\mathcal{A} * \mathcal{B}$  is again an element of  $\mathcal{M}(Q)$ . This defines a monoid structure on  $\mathcal{M}(Q)$  which is too unwieldy to study directly. Observe that we changed the order of the terms wrt. the definition given in [2]. That is, we will work with the *opposite* monoid of [2].

On the other hand, the *Reineke composition monoid* is very tractable. It is the submonoid  $\mathcal{C}(Q)$  of  $\mathcal{M}(Q)$  generated by the vertex-representation spaces  $R_i = \text{rep}_{\delta_i} Q$ . These generators satisfy specific commutation relations which can be read off from the quiver structure, see [2, §5]. For example, if there are no arrows between  $v_i$  and  $v_j$  then

$$R_i * R_j = R_j * R_i$$

and if there are no arrows from  $v_i$  to  $v_j$  but  $n$  arrows from  $v_j$  to  $v_i$ , then

$$\begin{cases} R_i^{*(n+1)} * R_j = R_i^{*n} * R_j * R_i \\ R_i * R_j^{*(n+1)} = R_j * R_i * R_j^{*n} \end{cases}$$

For more details on the structure of  $\mathcal{C}(Q)$  we refer to [2, §5].

There is a relation between  $\mathcal{C}(Q)$  and the left- and right- non-commutative topologies  $\Lambda_A^l$  and  $\Lambda_A^r$ . Because the Zariski topology on  $\text{simp } A$  is the discrete topology on the set  $\{S_1, \dots, S_k\}$  of vertex simplices, it is important to understand  $\mathcal{O}_w^r$  where  $w$  is a word in the  $S_i$ , say  $w = S_{i_1} S_{i_2} \dots S_{i_u}$ . In fact, we could have based our definition of a one-sided non-commutative topology on the set  $\mathcal{L}_A$  of *irreducible* open subsets of  $\text{simp } A$  and then these basic opens would be all. If  $\mathcal{C}$  is a  $GL(\alpha)$ -stable subset of  $\text{rep}_\alpha Q$  with  $|\alpha| = n$ , we will denote the subset  $GL_n \times^{GL(\alpha)} \mathcal{C}$  of  $\text{rep}_n A$  by  $\tilde{\mathcal{C}}$ .

### Proposition 1

$$\mathcal{O}_w^l = \bigcup_{w'} \tilde{\mathcal{A}}_{w'} \quad \text{resp.} \quad \mathcal{O}_w^r = \bigcup_{w'} \tilde{\mathcal{A}}_{w'}$$

where  $\mathcal{A}_{w'}$  is a  $*$ -word in the generators  $R_i$  of the composition monoid such that  $w'$  can be rewritten (using the relations in  $\mathcal{C}(Q)$ ) in the form

$$w' = R_{i_1} * R_{i_2} * \dots * R_{i_u} * w'' \quad \text{resp.} \quad w' = w'' * R_{i_1} * R_{i_2} * \dots * R_{i_u}$$

for another  $*$ -word  $w''$ .

Also, the equivalence relation introduced before can be expressed in terms of  $\mathcal{C}(Q)$ . If  $w = S_{i_1}S_{i_2}\dots S_{i_u}$  and  $w' = S_{j_1}S_{j_2}\dots S_{j_v}$  such that  $w \cup w' = \{S_{k_1}, \dots, S_{k_w}\}$ , then

**Proposition 2**  $\mathcal{O}_w^l \approx \mathcal{O}_{w'}^l$  if and only if every  $*$ -word  $v = R_{a_1} * \dots * R_{a_z}$  containing in it distinct factors  $R_{k_1}, \dots, R_{k_w}$  which can be brought in  $\mathcal{C}(Q)$  in the form

$$v = R_{i_1} * \dots * R_{i_u} * v'$$

can also be written in the form

$$v = R_{j_1} * \dots * R_{j_v} * v''$$

(and conversely). A similar result describes  $\mathcal{O}_w^r \approx \mathcal{O}_{w'}^r$ .

In particular, in this setting there will be hardly any *idempotent* basic opens (that is, satisfying  $\mathcal{O}_w^r \wedge \mathcal{O}_w^r \approx \mathcal{O}_w^r$ ). Clearly, if  $\{S_{e_1}, \dots, S_{e_a}\}$  are simples such that the quiver restricted to  $\{v_{e_1}, \dots, v_{e_a}\}$  has no arrows, then any word  $w$  in the  $S_{e_j}$  gives an idempotent  $\mathcal{O}_w^r$ . In the following section we will give an example where *every* basic open is idempotent and hence we get a commutative topology.

## 5 The commutative case.

If  $A$  is a commutative affine  $\mathbb{k}$ -algebra, then any simple representation is one-dimensional,  $\text{simp } A = X_A$  the affine (commutative) variety corresponding to  $A$  and the Zariski topologies on both sets coincide. Still, one can define the non-commutative topologies on  $\text{rep } A$ . However,

**Proposition 3** *If  $A$  is a commutative affine  $\mathbb{k}$ -algebra, then both  $\Lambda_A^l$  and  $\Lambda_A^r$  are commutative topologies. That is, for all words  $w$  and  $w'$  in  $\mathcal{L}_A$  we have*

$$\mathcal{O}_w^l \wedge \mathcal{O}_{w'}^l \approx \mathcal{O}_{w'}^l \wedge \mathcal{O}_w^l \quad \text{and} \quad \mathcal{O}_w^r \wedge \mathcal{O}_{w'}^r \approx \mathcal{O}_{w'}^r \wedge \mathcal{O}_w^r$$

**Proof.** We claim that every basic open  $\mathcal{O}_w^l$  is idempotent. Observe that all simple  $A$ -representations are one-dimensional and that there are only self-extensions of those, that is, if  $S$  and  $T$  are non-isomorphic simples, then  $\text{Ext}_A^1(S, T) = 0 = \text{Ext}_A^1(T, S)$ . However, there are self-extensions with the dimension of  $\text{Ext}_A^1(S, S)$  being equal to the dimension of the tangent space at  $X_A$  in the point corresponding to  $S$ . As a consequence we have for any Zariski open subsets  $U$  and  $V$  of  $X_A$  that

$$\mathcal{O}_{UV}^l = \mathcal{O}_{VU}^l$$

as we can change the order of the filtration factors (a representation  $M$  is the direct sum of submodules  $M_1 \oplus \dots \oplus M_s$  with each  $M_i$  concentrated in a single simple  $S_i$  and we can add the successive  $S_i$  factors of  $M$  at any wanted place in the filtration sequence). Hence, for every word  $w$  we have that

$$\mathcal{O}_w^l \approx \mathcal{O}_w^l \wedge \mathcal{O}_w^l$$



and also for any pair of words  $w$  and  $w'$  we have that

$$\mathcal{O}_w^l \wedge \mathcal{O}_{w'}^l = \mathcal{O}_{ww'}^l = \mathcal{O}_{w'w}^l = \mathcal{O}_{w'}^l \wedge \mathcal{O}_w^l$$

Observe that in [5] it is proved that a non-commutative topology in which every basic open is idempotent is commutative. We cannot use this here as the proof of that result uses both the left- and right- conditions. However, we are dealing here with a very simple example.  $\square$

## 6 Quotient varieties versus moduli spaces.

Having defined a one-sided non-commutative topology on  $\text{rep } A$  we can ask about the induced topology on the quotient variety  $\text{iss } A$  of all isomorphism classes of semi-simple  $A$ -representations or on the moduli space  $\text{moduli}_\theta A$  with respect to a certain stability structure  $\theta$ , cfr. [3]. Experience tells us that it is a lot easier to work with quotient varieties than with moduli spaces and non-commutative topology may give a partial explanation for this.

Indeed, as the points of  $\text{iss } A$  are semi-simple representations, it is clear that the induced non-commutative topology on  $\text{iss } A$  is in fact commutative. However, as the points of  $\text{moduli}_\theta A$  correspond to isomorphism classes of direct sums of stable representations (not simples!), the induced non-commutative topology on  $\text{moduli}_\theta A$  will in general remain non-commutative. Still, in nice examples, such as representations of quivers, one can define another non-commutative topology on  $\text{moduli}_\theta A$  which does become commutative. Use universal localization to cover  $\text{moduli}_\theta A$  by opens isomorphic to  $\text{iss } A_\Sigma$  for some families  $\Sigma$  of maps between projectives and equip  $\text{moduli}_\theta A$  with a non-commutative topology (which then will be commutative!) obtained by gluing the induced non-commutative topologies on the  $\text{rep } A_\Sigma$ .

## 7 Generalizations.

It should be evident that our construction can be carried out verbatim in the setting of any Artinian Abelian category (that is, an Abelian category having Jordan-Hölder sequences) as soon as we have a natural topology on the set of simple objects. In fact, the same procedure can be applied when we have a left (or right) non-commutative topology on the simples.

In fact, the construction may even be useful in Abelian categories in which every object is filtered by special objects on which we can define a (one-sided) (non-commutative) topology.

## References

- [1] Claudio Procesi, *Rings with polynomial identities*, Marcel Dekker (1973)

- [2] Markus Reineke, *The monoid of families of quiver representations*, Proc. London Math. Soc. **84** (2002) 663-685
- [3] Alexei Rudakov, *Stability for an Abelian category*, J. Alg. **197** (1997) 231-245
- [4] Fred Van Oystaeyen, *Prime spectra in noncommutative algebra*, Lect. Notes Math. 444, Springer (1975)
- [5] Fred Van Oystaeyen, *Algebraic geometry for associative algebras*, Marcel Dekker (2000)
- [6] Fred Van Oystaeyen, *Virtual topology and functor geometry*, monograph, to appear.